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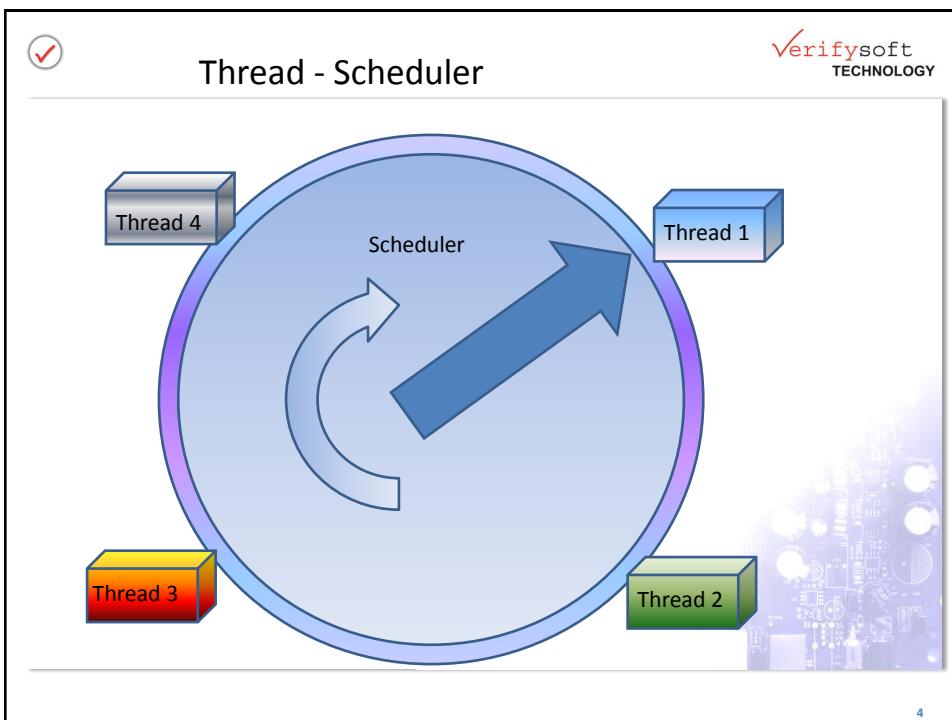
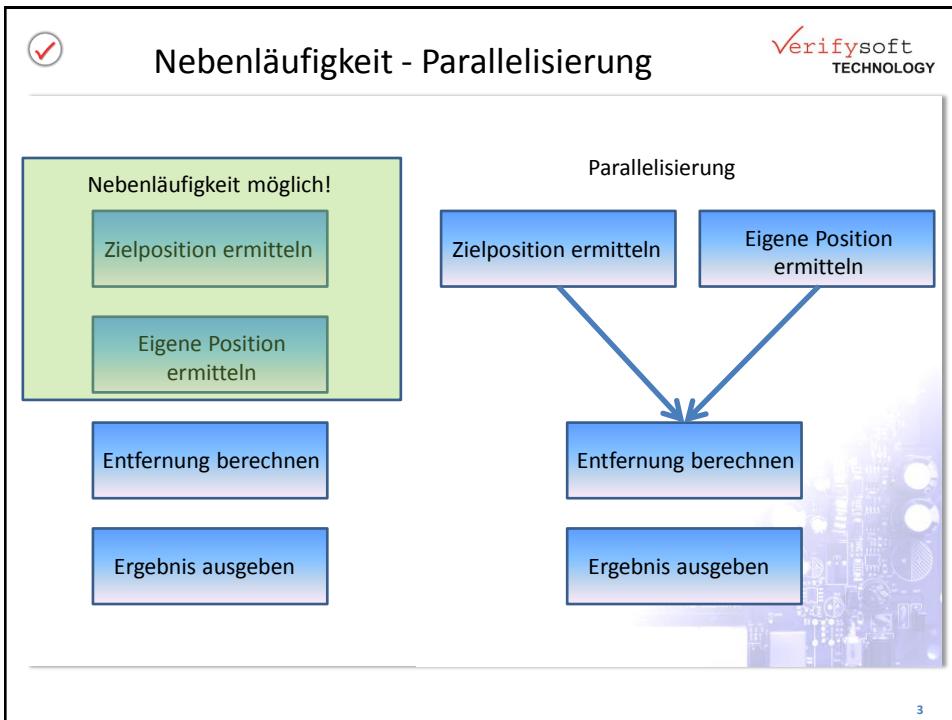
# Auffinden von Nebenläufigkeitsfehlern durch statische Codeanalyse

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## Agenda

- Was ist Nebenläufigkeit?
- Was ist Parallelisierung?
- Zuweisung von CPU-Zeit durch den Scheduler
- Motivation zur Implementierung von Nebenläufigkeit?
- Was sind Nebenläufigkeitsfehler?
- Warum sind Nebenläufigkeitsfehler so gefürchtet?
- Beispiel Data Race
- Beispiel Dead Lock (Nested Lock)
- Wie kann die statische Codeanalyse helfen?

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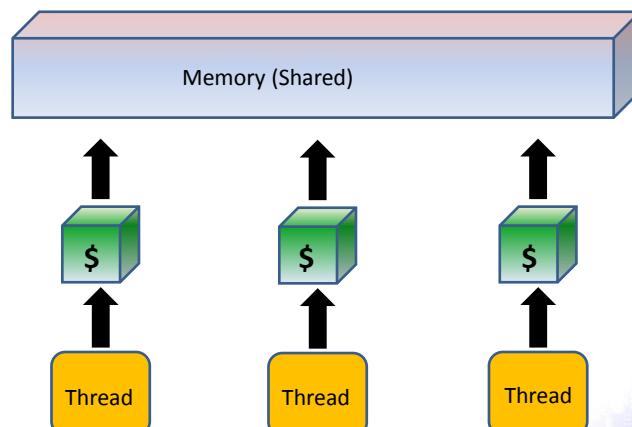
Motivation zur Implementierung von Nebenläufigkeit 

### Multicore Mikrocontroller (Auswahl)

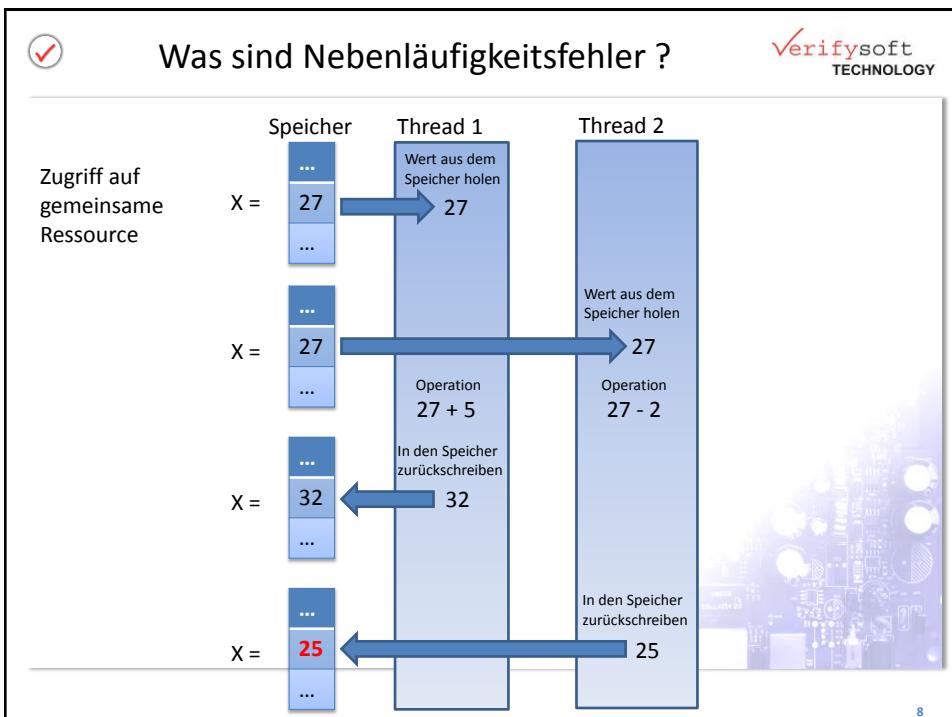
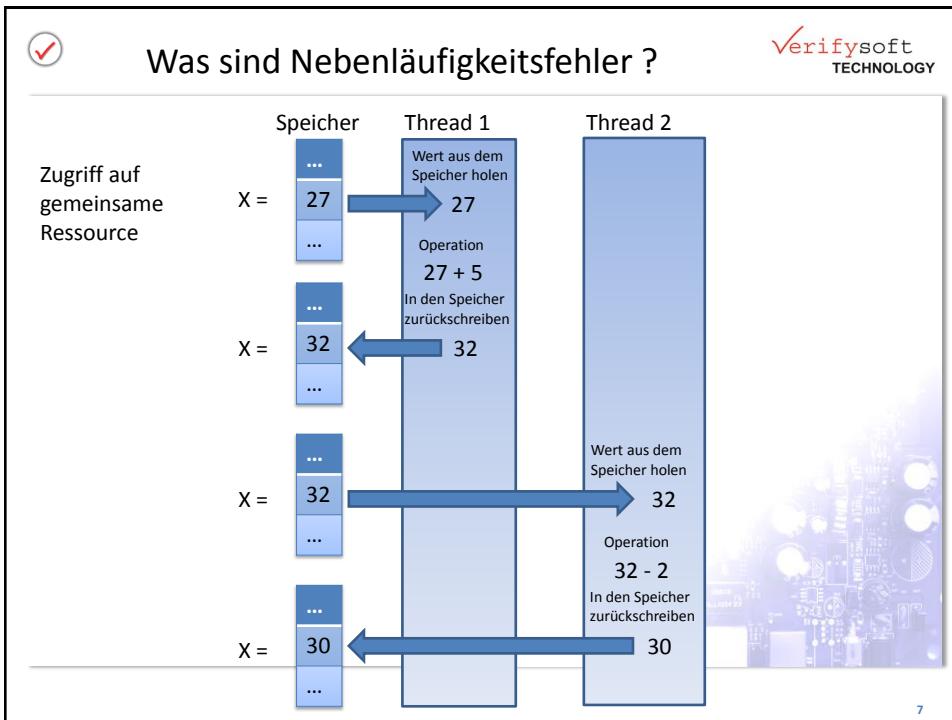
Hersteller	Typ	Cores	Bitness
XMOS	L-Serie	4 - 16	64
Freescale	MPC5777M MCU	4	32
Infinion	AURIX™ Family – TC29xT	3	32
STMicroelectronics	SPC56EL70	2	32
Parallax	Propeller	8	32

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Multithreaded Programming Model 



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Warum sind Nebenläufigkeitsfehler so gefürchtet ?

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### Problematik des Auffindens zur Laufzeit

- In der Regel kein deterministisches Auftreten
- Verhalten oft lastabhängig
- Einsatz von Prüfwerkzeugen (z. B. Debugger) ändert das Laufzeitverhalten oft derart, dass der Fehler nicht mehr auftritt

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Heisenbergsche Unschärferelation

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$$\Delta x * \Delta p \sim h$$

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## Beispiel: Data Race

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```
int main(void){
    DWORD    threadId_1, threadId_2;
    HANDLE   threadHandle_1, threadHandle_2;
    unsigned long segment_info0[3] = { 0, 5000000, 10000000 };
    unsigned long segment_info1[3] = { 5000001, 10000000, 10000000 };

    threadHandle_1 = CreateThread(NULL, 0, &calculate, (LPVOID)segment_info0, 0, &threadId_1);
    if (threadHandle_1 == NULL) {
        exit(EXIT_FAILURE);
    }
    threadHandle_2 = CreateThread(NULL, 0, &calculate, (LPVOID)segment_info1, 0, &threadId_2);
    if (threadHandle_2 == NULL) {
        exit(EXIT_FAILURE);
    }

    WaitForSingleObject(threadHandle_1, INFINITE);
    WaitForSingleObject(threadHandle_2, INFINITE);

    CloseHandle(threadHandle_1);
    CloseHandle(threadHandle_2);

    printf("PI = %25.20f\n", sum / segment_info0[2]);

    return 0;
}
```



## Fehlende Synchronisierung -> Data Race

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```
double sum;

DWORD WINAPI calculate(LPVOID params){
    double d, w;
    unsigned long l;
    unsigned long *part;

    part = (unsigned long*)params;
    d = 1.0 / part[2];
    for (l = part[0]; l < part[1]; ++l){
        w = d * (l + 0.5);
        sum += 4.0 / (1.0 + w * w);
    }
    return 0;
}
```

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## Data Race – Lösung: Synchronisierung

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```
double sum;
CRITICAL_SECTION lock;

DWORD WINAPI calculate(LPVOID params){
    double d, w;
    unsigned long l;
    unsigned long *part;

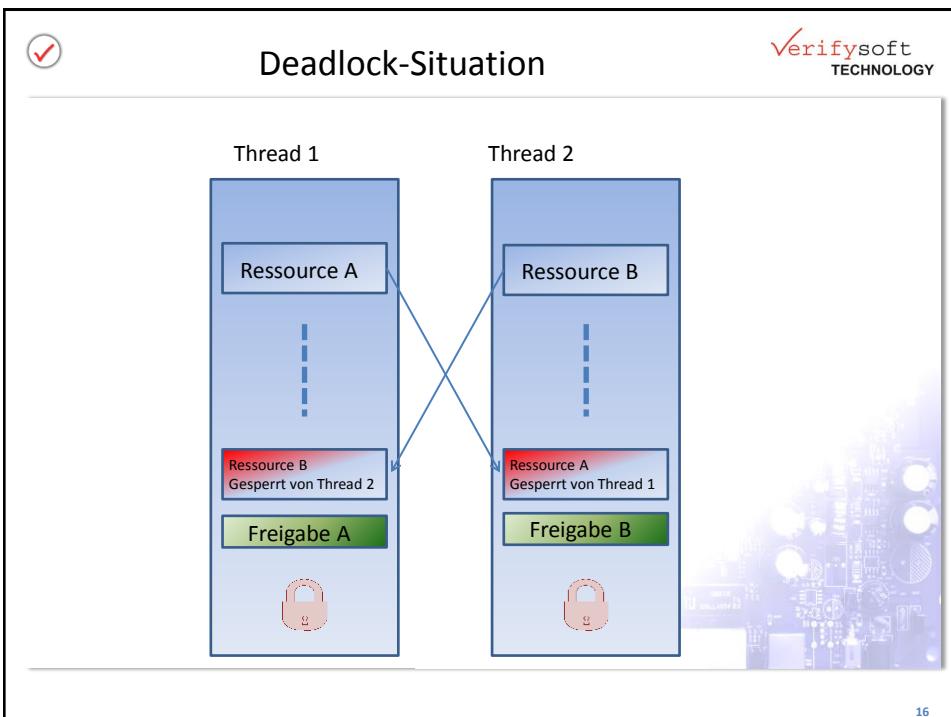
    InitializeCriticalSection(&lock);
    part = (unsigned long*)params;
    d = 1.0 / part[2];
    for (l = part[0]; l < part[1]; ++l){
        w = d * (l + 0.5);
        EnterCriticalSection(&lock);
        sum += 4.0 / (1.0 + w * w);
        LeaveCriticalSection(&lock);
    }
    DeleteCriticalSection(&lock);
    return 0;
}
```

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Aufdeckung durch statische Codeanalyse

CODESONAR Search code in this analysis for [Search] | Advanced Search  
 Home > DataRace > DataRace analysis 1 > Warning 161.3391  
**Data Race** at datarace\_windows.cpp:16 No properties have been set | edit properties  
 Jump to warning location | warning details...  
 Show Events | Change View | Options  
**thread 1**  
**calculate** (c:\tests\datarace\_windows\datarace\_windows\datarace\_windows.cpp)  
 7 WORD WINAPI calculate(LVOID params){  
 8 double d, w;  
 9 unsigned long i;  
 10 unsigned long \*part;  
 11  
 12 part = (unsigned long\*)params;  
 13 d = 1.0 / part[2];  
 14 for (i = part[0]; i < part[1]; ++i){  
 15 r = (i + 0.5) \* d;  
 16 sum += 4.0 / (1.0 + w \* r);  
 Data Race  
 This code reads from global variable sum.  
 • The other thread writes to sum. See [other access](#).  
 • No locks are currently held so a race with the other thread may occur.  
 • Compilers and processors reorder accesses to shared variables, so even source code that looks safe can be vulnerable to data races.  
 The issue can occur if the highlighted code executes.  
 Show: All events | Only primary events

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 Typische Deadlock-Konstellation

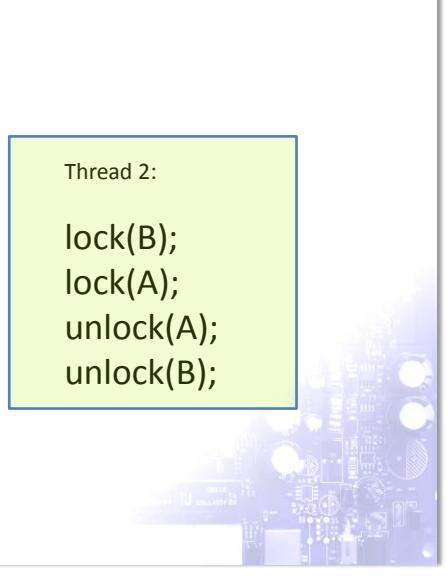
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Thread 1:

```
lock(A);
lock(B);
unlock(B);
unlock(A);
```

Thread 2:

```
lock(B);
lock(A);
unlock(A);
unlock(B);
```



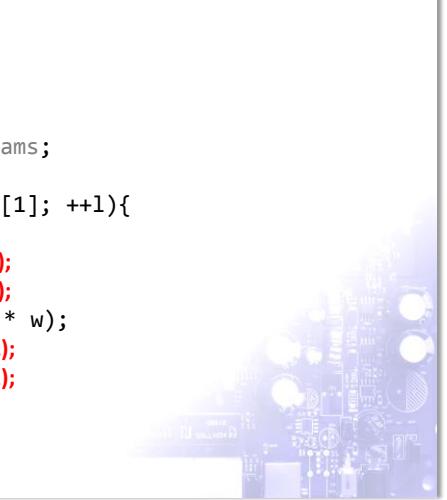
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 Nested Lock

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```
DWORD WINAPI calculate_1(LPVOID params){
    double d, w;
    unsigned long l;
    unsigned long *part;

    InitializeCriticalSection(&lock_1);
    InitializeCriticalSection(&lock_2);
    part = (unsigned long*)params;
    d = 1.0 / part[2];
    for (l = part[0]; l < part[1]; ++l){
        w = d * (1 + 0.5);
        EnterCriticalSection(&lock_1);
        EnterCriticalSection(&lock_2);
        sum += 4.0 / (1.0 + w * w);
        LeaveCriticalSection(&lock_2);
        LeaveCriticalSection(&lock_1);
    }
    DeleteCriticalSection(&lock_1);
    DeleteCriticalSection(&lock_2);
    return 0;
}
```



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 Nested Lock 

```

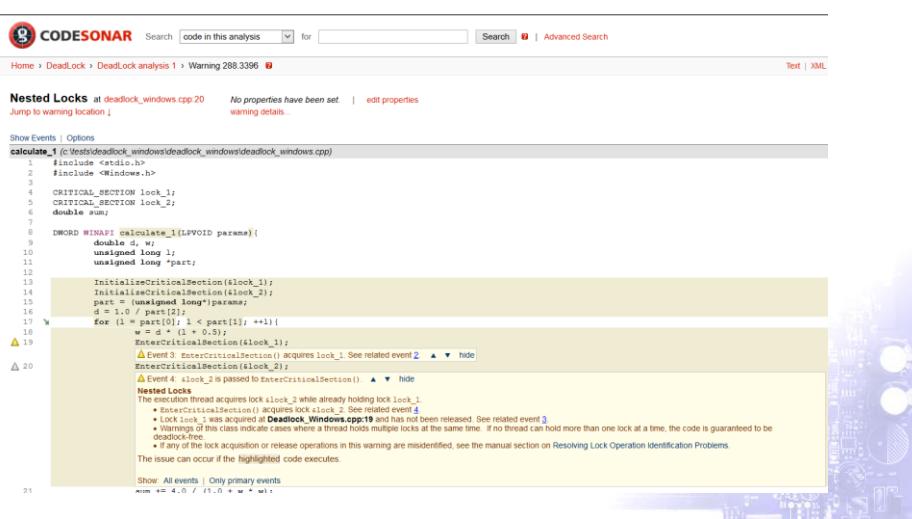
DWORD WINAPI calculate_2(LPVOID params){
    double d, w;
    unsigned long l;
    unsigned long *part;

    InitializeCriticalSection(&lock_1);
    InitializeCriticalSection(&lock_2);
    part = (unsigned long*)params;
    d = 1.0 / part[2];
    for (l = part[0]; l < part[1]; ++l){
        w = d * (1 + 0.5);
        EnterCriticalSection(&lock_2);
        EnterCriticalSection(&lock_1);
        sum += 4.0 / (1.0 + w * w);
        LeaveCriticalSection(&lock_1);
        LeaveCriticalSection(&lock_2);
    }
    DeleteCriticalSection(&lock_1);
    DeleteCriticalSection(&lock_2);
    return 0;
}

```

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 Aufdeckung durch statische Codeanalyse 



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